Design, development, and analysis of a compressible fluid dynamics solver capable of exceeding one quadrillion degrees of freedom

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Collaborative effort



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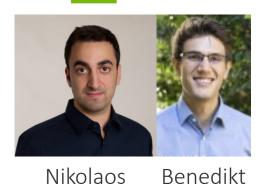
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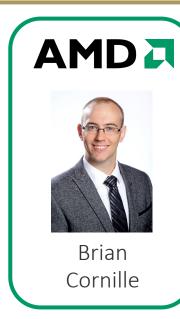




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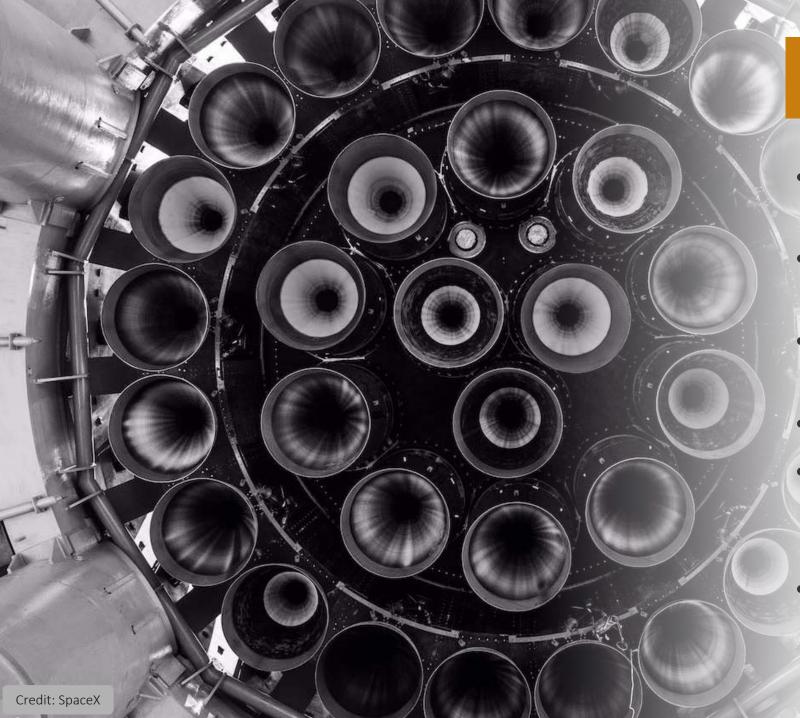
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Introduction

- Large arrays of small rockets are the new standard
- Interactions between nearby rockets produce new problems
- Simulation resolution is limited by available GPU memory
- More resolution = more rockets
- Existing tools for shock-laden flows are poorly conditioned at lower precision
- Lower precision and linear numerics are faster

Timeline of biggest CFD simulations

- 10T grid cells on BlueGene/Q
- Enabled by susbstanial CPU memory
- Nonlinear state-of-the-art numerics
- Time to solution ~ 1 month for meaningful time scales

2013^[1] (GB Winner)

- 10T grid cells on OLCF Frontier
- Nonlinear state-of-the-art numerics
- Time to solution: hours to days
- Compressible

Oct. 2024^[3]

- 200T cells on OLCF Frontier
- 113T cells on LLNL El Cap

Today

- 100T cells on JSC Jupiter
- Compressible flow solver
- Novel numerics
- Time to solution: hours to days

 35T grid cells on OLCF Frontier

Sep. 2024^[2]

- Incompressible flow
- Time-to-solution dominated by all-to-all communication
- Time to solution: hours to days

- [1] Rossinelli et al. (2013) Proceedings of SC '13
- [2] Yeung et al. (2025) CPC
- [3] Sathyanarayana et al. (2025) JPDC

Summary of contributions

- Information geometric regularization foregoes nonlinear viscous shock capturing, enabling linear off-the-shelf numerical schemes and sequential summation of right-hand side contributions.
- Unified addressing on tightly coupled CPU-GPU and APU platforms increases total problem size with negligible performance hit.
- FP32 compute and FP16 storage further reduce memory use while remaining numerically stable, enabled by the algorithm's well-conditioned numerics.
- Reduce memory footprint 25-fold over state-of-the-art. Improve time and energy-to-solution factors of 4 and 5.4, compared to an optimized implementation of state-of-the-art methods.
- First CFD simulation exceeding 200T grid points and 1 quadrillion degrees of freedom, improving on previous largest simulations by a factor of 20.

How big is a quadrillion???

1 quadrillion = $1,000,000,000,000,000 = 1 \times 10^{15}$

1 quadrillion seconds $\approx 31.7 \ million$ years



- Model and numerical method
- Basic implementation details
- System-specific design
- Performance and results

Model and numerical method

Navier-Stokes equations for a single fluid

$$\frac{\partial \rho}{\partial t} + \nabla \cdot (\rho \mathbf{u}) = 0,$$

$$\frac{\partial(\rho \boldsymbol{u})}{\partial t} + \nabla \cdot (\rho \boldsymbol{u} \otimes \boldsymbol{u} + p\mathbf{I} - \mathbf{T}) = \mathbf{0},$$
$$\frac{\partial E}{\partial t} + \nabla \cdot [(E + p)\boldsymbol{u} - \boldsymbol{u} \cdot \mathbf{T}] = 0,$$

Old solution method:

- High-order finite volume solver
- HLLC Riemann solver
- WENO spatial reconstructions
- Requires converting between conservative and primitive variables for stability
- Expensive, nonlinear, and illconditioned at lower floating-point precisions

Model and numerical method

Navier-Stokes equations for a single fluid

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$$\frac{\partial(\rho \boldsymbol{u})}{\partial t} + \nabla \cdot (\rho \boldsymbol{u} \otimes \boldsymbol{u} + p\mathbf{I} - \mathbf{T}) = \mathbf{0},$$

$$\frac{\partial E}{\partial t} + \nabla \cdot [(E + p)\boldsymbol{u} - \boldsymbol{u} \cdot \mathbf{T}] = 0,$$

New solution method (IGR):

- Add some terms to the equation
- Solve using linear numerics and in purely conservative variables

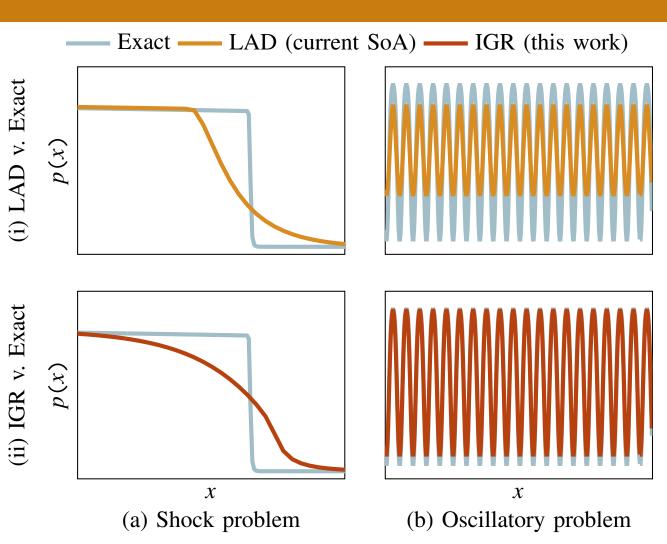
$$\frac{\partial \rho}{\partial t} + \nabla \cdot (\rho \boldsymbol{u}) = 0,$$

$$\frac{\partial (\rho \boldsymbol{u})}{\partial t} + \nabla \cdot (\rho \boldsymbol{u} \otimes \boldsymbol{u} + (p + \boldsymbol{\Sigma})\mathbf{I} - \mathbf{T}) = \mathbf{0},$$

$$\frac{\partial E}{\partial t} + \nabla \cdot [(E + p + \boldsymbol{\Sigma})\boldsymbol{u} - \boldsymbol{u} \cdot \mathbf{T}] = 0,$$

$$\alpha \left[\operatorname{tr} (\nabla \boldsymbol{u})^2 + \operatorname{tr}^2 (\nabla \boldsymbol{u}) \right] = \frac{\Sigma}{\rho} - \alpha \nabla \cdot \left(\frac{\nabla \Sigma}{\rho} \right)$$

Benefits of IGR



- Existing approach (LAD) yields solutions that are not smooth to higher orders
 - Can lead to spurious oscillations at discontinuities and dissipation of oscillatory features
- IGR replaces shocks with highorder smooth profiles to reduce oscillations near shocks and dissipation of oscillatory features

Credit: Cao and Schafer, 2024

Solution overview

$$q^{(2)}=q^{(1)},$$
 $q^{(1)}=q^{(1)}+\Delta t rac{\partial q^{(1)}}{\partial t},$ $q^{(1)}=rac{3}{4}q^{(2)}+rac{1}{4}q^{(1)}+rac{1}{4}\Delta t rac{\partial q^{(1)}}{\partial t},$ $q^{(1)}=rac{3}{4}q^{(2)}+rac{1}{4}q^{(1)}+rac{1}{4}\Delta t rac{\partial q^{(1)}}{\partial t},$ $q^{(1)}=rac{1}{3}q^{(2)}+rac{2}{3}q^{(1)}+rac{2}{3}\Delta t rac{\partial q^{(1)}}{\partial t}$

RHS Calculation

$$\frac{\partial \boldsymbol{q}}{\partial t} = \begin{cases} \frac{\partial \rho}{\partial t} = -\nabla \cdot (\rho \boldsymbol{u}), \\ \frac{\partial (\rho \boldsymbol{u})}{\partial t} = -\nabla \cdot (\rho \boldsymbol{u} \otimes \boldsymbol{u} + (p + \Sigma)\mathbf{I} - \mathbf{T}, \\ \frac{\partial E}{\partial t} = -\nabla \cdot [(E + p + \Sigma)\boldsymbol{u} - \boldsymbol{u} \cdot \mathbf{T}], \end{cases}$$

Elliptic solve for entropic pressure

$$\frac{\Sigma}{\rho} - \alpha \nabla \cdot \left(\frac{\nabla \Sigma}{\rho}\right) = \underbrace{\alpha \left[\operatorname{tr} \left(\nabla \boldsymbol{u}\right)^{2} + \operatorname{tr}^{2} \left(\nabla \boldsymbol{u}\right) \right]}_{\Sigma_{\text{rhs}}}$$

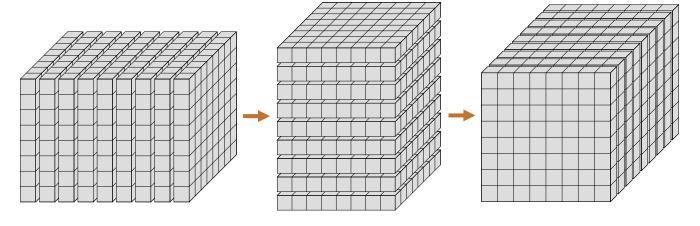


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Domain decomposition and communication

~0.4% of domain

Block decomposition communication pattern



Example max problem size decompositions

$N_{ m proc}$	Proc Disc.	Total (T)
128	$4 \times 4 \times 8$	0.340
384	$8 \times 8 \times 6$	1.02
1024	$16 \times 16 \times 8$	2.73
3072	$16 \times 16 \times 12$	8.18
8192	$32 \times 32 \times 16$	21.8
75264	$48 \times 49 \times 32$	200

Storage and I/O

```
restart_data
      lusture_0.dat
      lusture_100.dat
      lusture_0.dat
      lusture_100.dat
```

Pros:

- Simple
- Fast

Cons:

- O(10⁶) files for leadership scale simulation
- Lots of concurrent metadata creation

$$(1 \times 10^{15} \text{ scalars}) \times \left(16 \frac{\text{bits}}{\text{scalar}}\right) / \left(8 \times 10^{12} \frac{\text{bits}}{\text{terabyte}}\right) = 2000 \text{ terabytes} = 2 \, \textbf{Petabytes}$$

GPU programming landscape

	OpenMP			OpenACC		
Compiler	NV GPUs	AMD GPUs	FP16 Atomics	NV GPUs	AMD GPUs	FP16 Atomics
AMD	×			×	×	×
CCE			×			×
NVHPC		X	?		X	*

- Modern supercomputers are procured from two hardware vendors, so portability is important!
- OpenMP is generally better supported, but OpenACC is generally faster when the necessary features are available
- If we want the best performance we can get on all hardware we need to support everything

Portable GPU offload using directive based programming and macros

Source code with macros

Preprocessor

- OpenMP and OpenACC are generated from one version of source code
- Developers don't need in depth understanding of directive tools

```
!$omp target teams distribute parallel do simmd &
!$omp defaultmap(firstprivate:scalar) &
!$omp defaultmap(tofrom:aggregate) &
!$omp defaultmap(present:allocatable) &
!$omp defaultmap(present:pointer) &
!$omp collapse(3) private(...)
do l = 0, Nz     ! Z-direction
                                                OpenMP
   do k = 0, Ny ! Y-direction
                                                  Code
       do j = 0, Nx ! X-direction
       !!> Core kernel,
           !!> O(1000) arithmetic operations
       end do
   end do
end do
!$omp end target teams distribute parallel do
```

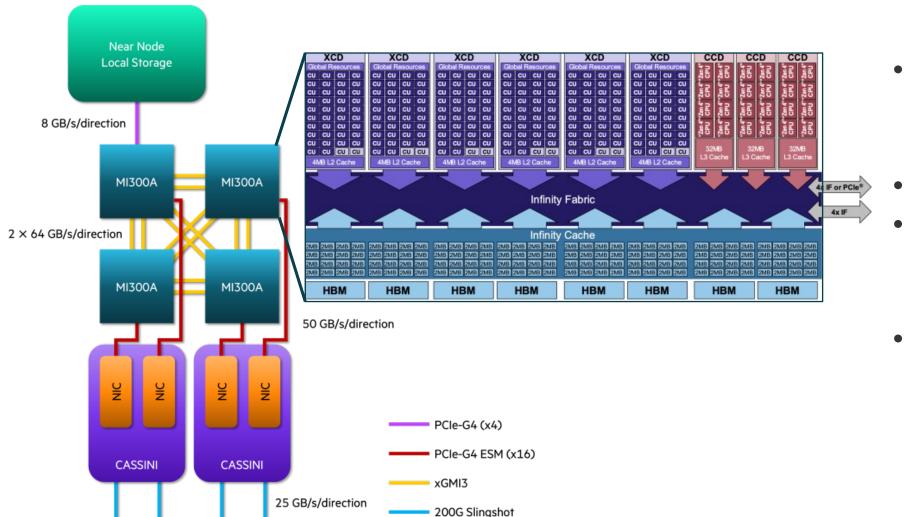
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System Summary

	Node Configuration	# Nodes	Memory [Node, System]	Peak Power	Rmax	TOP500
LLNL El Capitan	4 AMD MI300A APU	11136	$[512\mathrm{GB},5.6\mathrm{PB}]\;\mathrm{APU}$	$34.8\mathrm{MW}$	$1742\mathrm{PFLOPs}$	1
OLCF Frontier	4 AMD MI250X GPU 1 AMD Trento CPU	9472	$[512\mathrm{GB},4.8\mathrm{PB}]\mathrm{GPU}$ $[512\mathrm{GB},4.8\mathrm{PB}]\mathrm{CPU}$	$24.6\mathrm{MW}$	1353 PFLOPs	2
CSCS Alps	4 NVIDIA GH200 (4 Grace CPU, 4 Hopper GPU)	2688	[384 GB, 1.0 PB] GPU [480 GB, 1.3 PB] CPU	$7.1\mathrm{MW}$	$435\mathrm{PFLOPs}$	8

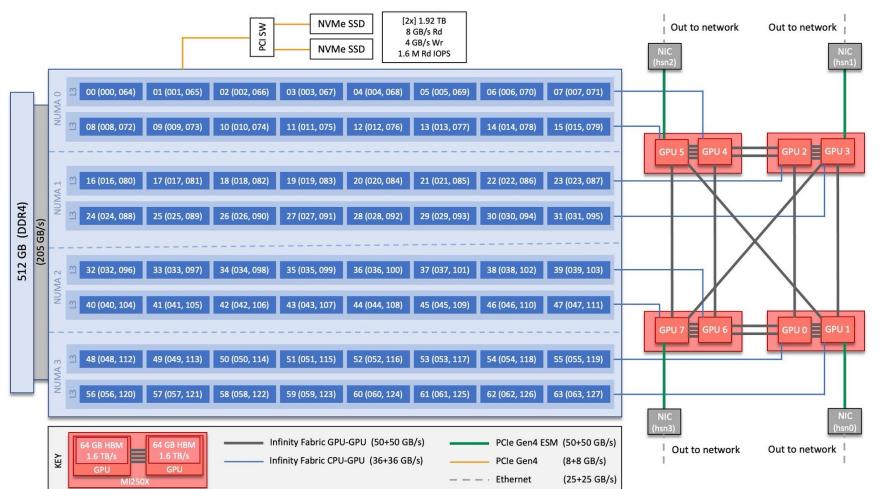
- All systems of interest are within the top 10 fastest computers measured by the HPL benchmark
- CSCS Alps is "little brother" to 6k node JSC Jupiter which is ranked 3rd now
- >95% of total system memory is used on El Capitan and Frontier
- >85% of total system memory is used on Alps
- Shared critical features:
 - Ability to leverage unified address space between CPU and GPU
 - High performance networking and leadership class scale

Tightly coupled GPU/APU architectures (AMD MI300A, LLNL El Capitan)



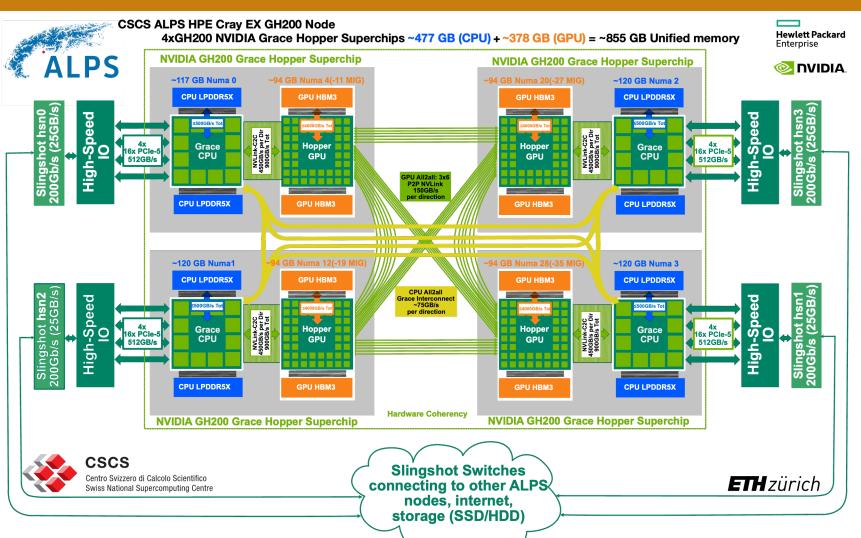
- 24 Zen4 CPU cores bonded to 6 AMD XCD chiplets
- 4 APUs per node
- Memory is universally addressable in address space and physical space
- Zero host-to-device transfers to perform because the host and device are one

Tightly coupled GPU/APU architectures (AMD MI250X, OLCF Frontier)



- 4 Mi250X GPUs (8 GCDs)
 connected to one CPU via
 36+36 GB/s links
- 1 CPU and 4 GPUs per node
- Each chip has its own memory
- Unified address space made possible through careful memory allocation and infinity fabric

Tightly coupled GPU/APU architectures (NVIDIA GH200, CSCS Alps)



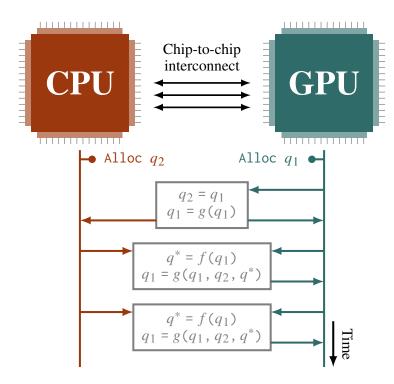
- Grace CPU and Hopper GPU connected via 900 GB/s chip-to-chip (C2C) link
- 4 CPUs and 4 GPUs per node
- Each chip has its own memory
- Unified address space made possible via highspeed interconnect, though memory regions are physically separate

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Unified memory strategy

MI250x - Allocate q₂ on the using hipMallocManaged and advise runtime *not* to make device copy

GH200 – Allocate q₂ and use cudaMemAdvise to keep memory on the host



MI250x - Allocate q₁ on the GPU using hipMalloc and advise runtime not to make a host copy

GH200 – Allocate q₁ and pin memory with cudaMemAdvise

• GH200's 900 GB/s link allows for storage of Σ and Σ_{rhs} on the host as well, further increasing problem size

- Model and numerical method
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- Performance and results

Performance: Grindtime

Grindtime

Device	Baseline (in-core)	IGR (in-core)	IGR (unified)	
GH200	16.89	3.83	4.18	FP64
MI250X GCD	69.72	13.01	19.81	
MI300A	29.50	†	7.21	
GH200	*N/A	2.70	2.81	FP32
MI250X GCD	*N/A	9.12	13.03	
MI300A	*N/A	†	4.19	
GH200 MI250X GCD MI300A	*N/A *N/A *N/A	3.06 22.63 †	3.07 24.71 17.39	FP16/32

^{*}Numerically unstable; †MI300A is always unified

Normalized by nanoseconds/grid cell/time step

- Baseline numerics unstable in lower precision
- AMD mixed precision slow due to using AMD's beta compiler and a pre-release of NVHPC SKD 25.9
- ~9% overhead in double precision on NVIDIA due to compiler regression
- <5% slowdown in single and mixed precision on NVIDIA thanks to 900 GB/s link
- 40-50% overhead in double and single precision on MI250x due to slower 36+36 GB/s link

Up to 6x reduction in time to solution

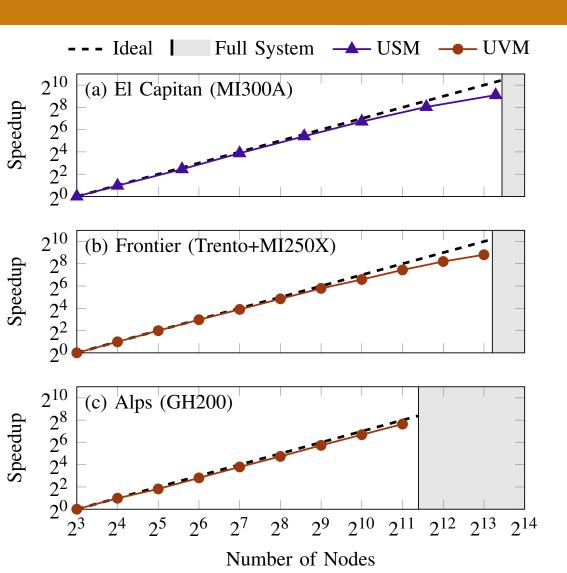
Performance: Power

Power

Energy (μJ)	El Capitan (MI300A)	Frontier (MI250X)	Alps (GH200)
Baseline IGR	$15.24 \\ 3.493$	10.67 1.982	$9.349 \\ 2.466$

- Calculated by sampling nvidia-smi and rocm-smi to get a steady state wattage and then multiply by time
- Measured in in double precision for an apples-to-apples comparison
- Findings show power consumption is proportional to runtime as a first-order approximation
- GH200 uses more power with novel numerics than current state-of-the-art, but algorithm is faster

Performance: Strong scaling



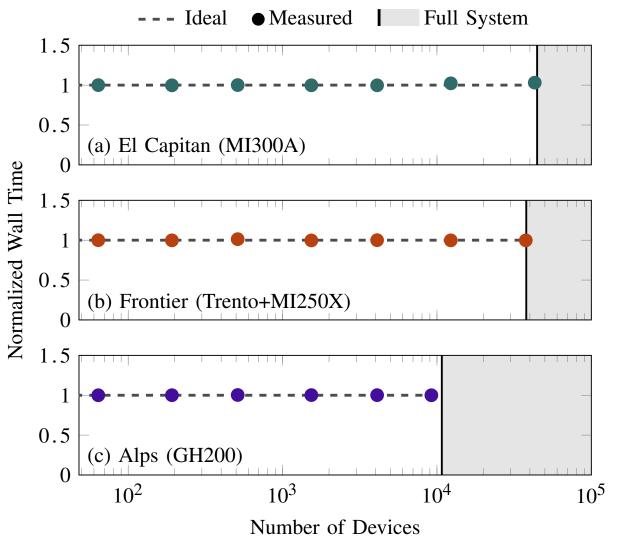
- 8-node base case does internode communication in all three physical dimensions
- Higher efficiency results in shorter time to solution for a given problem size

$$Efficiency = \frac{Actual\ Speedup}{Ideal\ Speedup}$$

Strong Scaling Efficiencies

Systemn	Base	32× Efficieny	Full System	
El Capitan	84 B	90%	10000 Nodes	44%
Frontier	170 B	90%	8192 Nodes	45%
Alps	134 B	86%	2048 Nodes	78%

Performance: Weak scaling



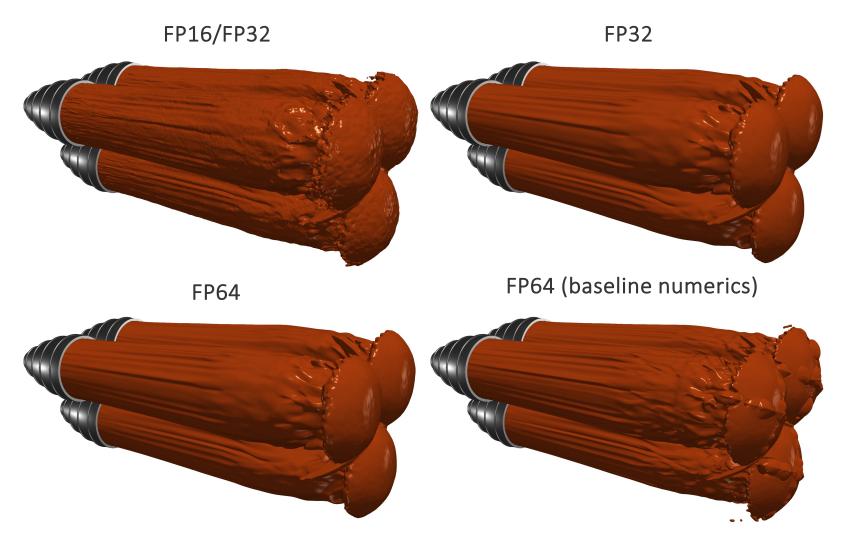
- 16-node base case does internode communication in all three physical dimensions (and makes for nice spacing)
- Efficiencies are all approx. 100%

Efficiency =
$$\frac{T_{\text{base case}}}{T}$$

Weak Scaling Efficiencies

System	Base	Max	Increase	Efficiency
El Capitan	168 B	113 T	$672 \times$	97%
Frontier	$341\mathrm{B}$		$588 \times$	99%
Alps	$268\mathrm{B}$	$38\mathrm{T}$	$143 \times$	99%

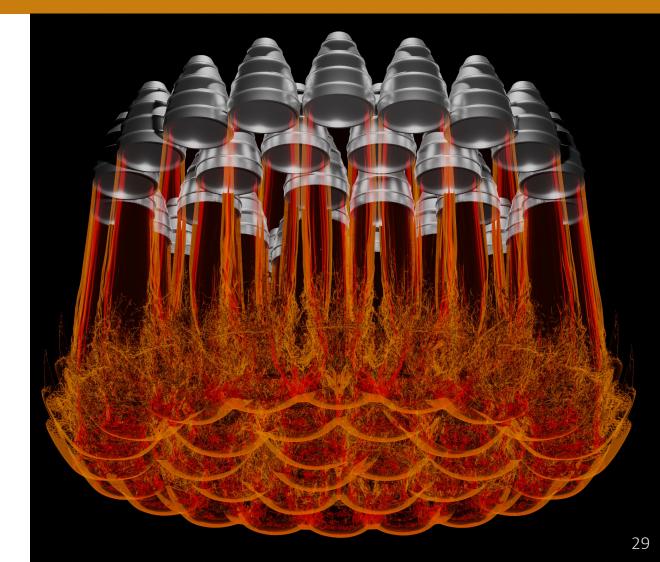
Results: Precision comparison



- FP64 and FP32 visually indistinguishable
- FP16/FP32 is visually different, though the main flow features are captured
- Numerical error leads to early onset of hydrodynamic instabilities
- FP64 baseline has grid aligned artifacts due to gird aligned shock capturing

Results: Super heavy booster configuration

- 3.3 trillion grid cells
- 16.5 trillion degrees of freedom
- 16 hours on 9200 GH200 GPUs



Lessons learned

- NVIDIA compilers are great, but they let you do a lot of things that aren't in the standards, so supporting new compilers can be challenging
- Compiler support for directive-based programming with lower floating-point precisions is still developing
- Heterogenous architectures are cool and allow you to do some interesting things
- Doing something no one has ever done before is difficult and time consuming





Questions?

Acknowledgements













